

INSTRUCTOR'S MANUAL
to accompany
Sixth Edition

**Fundamentals of
Digital Logic
and
Microcontrollers**

Sixth Edition

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CHAPTER 2

2.1 (a) $128 \ 64 \ 32 \ 16 \ 8 \ 4 \ 2 \ 1 \leftarrow \text{weighting}$
 $0 \ 1 \ 1 \ 1 \ 0 \ 1 \ 0 \ 1_2 = 64+32+16+4+1 = 117_{10}$

(b) $8 \ 4 \ 2 \ 1 \ .5 \ .25 \ .125 \leftarrow \text{weighting}$
 $1 \ 1 \ 0 \ 1 \ . \ 1 \ 0 \ 1_2 = 8+4+1+.5+.125 = 13.625_{10}$

(c) $8 \ 4 \ 2 \ 1 \ .5 \ .25 \ .125 \leftarrow \text{weighting}$
 $1 \ 0 \ 0 \ 0 \ . \ 1 \ 1 \ 1_2 = 8+.5+.25+.125 = 8.875_{10}$

2.2(a)

		Quotient	+	Remainder
152/2	=	76	+	0
76/2	=	38	+	0
38/2	=	19	+	0
19/2	=	9	+	1
9/2	=	4	+	1
4/2	=	2	+	0
2/2	=	1	+	0
1/2	=	0	+	1
152 ₁₀	=	1001 1000 ₂		

2.2(b)

		Quotient	+	Remainder
343/2	=	171	+	1
171/2	=	85	+	1
85/2	=	42	+	1
42/2	=	21	+	0
21/2	=	10	+	1
10/2	=	5	+	0
5/2	=	2	+	1
2/2	=	1	+	0
1/2	=	0	+	1
343 ₁₀	=	101010111 ₂		

2.3(a)

		Quotient	+	Remainder
1843/8	=	230	+	3
230/8	=	28	+	6
28/8	=	3	+	4
3/8	=	0	+	3
1843 ₁₀	=	3463 ₈		

2.3(b)

		Quotient	+	Remainder
1766/8	=	220	+	6
220/8	=	27	+	4
27/8	=	3	+	3
3/8	=	0	+	3
1766 ₁₀	=	3346 ₈		

2.4(a)

		Quotient	+	Remainder
1987/16	=	124	+	3
124/16	=	7	+	C
7/16	=	0	+	7
1987 ₁₀	=	7C3 ₁₆		

2.4 (b)

	Quotient	+	Remainder
3072/16 =	192	+	0
192/16 =	12	+	0
12/16 =	0	+	C
3072 ₁₀ =	C00 ₁₆		

2.5(a)

↓ Add leading zeros

$$\underbrace{001}_{\downarrow} \underbrace{101}_{\downarrow} \underbrace{011}_{\downarrow} \underbrace{100}_{\downarrow} \underbrace{101}_{\downarrow} = 15345_8$$

↓ Add 3 leading zeros

$$\underbrace{0001}_{\downarrow} \underbrace{1010}_{\downarrow} \underbrace{1110}_{\downarrow} \underbrace{0101}_{\downarrow} = 1AE5_{16}$$

(b)

↓ Add leading zero

$$\underbrace{011}_{\downarrow} \underbrace{000}_{\downarrow} \underbrace{011}_{\downarrow} \underbrace{100}_{\downarrow} \underbrace{110}_{\downarrow} \underbrace{000}_{\downarrow} \underbrace{011}_{\downarrow} = 3034603_8$$

$$\underbrace{1100}_{\downarrow} \underbrace{0011}_{\downarrow} \underbrace{1001}_{\downarrow} \underbrace{1000}_{\downarrow} \underbrace{0011}_{\downarrow} = C3983_{16}$$

2.6(a) Signed-magnitude form

48 in 7-bit can be represented as: $48 = 0110000_2$
 Therefore, $-48 = 10110000$ in sign magnitude form.
 52 in 7-bit = 0110100_2
 Hence, $+52 = 00110100_2$ in sign magnitude form

2.6(b) Ones complement form

$+48 = 0011\ 0000_2$
 $-48 = 1100\ 1111_2$
 $+52 = 0011\ 0100_2$

2.6(c) Two's complement form

$-48 = 1101\ 0000_2$
 $+52 = 0011\ 0100_2$

2.7

$1100\ 1100_2$ is an even number since the least significant bit is 0.
 $0010\ 0100_2$ is an even number since the least significant bit is 0.
 $0111\ 1001_2$ is an odd number since the least significant bit is 1.

2.8

532.372₁₀

	Quotient	+	Remainder
532/2 =	266	+	0
266/2 =	133	+	0
133/2 =	66	+	1
66/2 =	33	+	0
33/2 =	16	+	1
16/2 =	8	+	0
8/2 =	4	+	0
4/2 =	2	+	0
2/2 =	1	+	0
1/2 =	0	+	1

$$\begin{array}{r}
 \text{Hence, } 532_{10} = 10\ 0001\ 0100_2 \\
 \begin{array}{r}
 0.372 \\
 \underline{\times 2} \\
 0.744 \\
 \downarrow \\
 0
 \end{array}
 \quad
 \begin{array}{r}
 0.744 \\
 \underline{\times 2} \\
 1.488 \\
 \downarrow \\
 1
 \end{array}
 \quad
 \begin{array}{r}
 0.488 \\
 \underline{\times 2} \\
 0.976 \approx 1.000 \\
 \downarrow \\
 1
 \end{array} \\
 0.372_{10} \cong .011_2 \\
 532.372_{10} \cong 100\ 001\ 0100.011_2
 \end{array}$$

2.9 $15\text{FD}_{16} = 0001\ 0101\ 1111\ 1101_2$
 $26\text{EA}_{16} = 0010\ 0110\ 1110\ 1010_2$

2.10(a) $11264 = 0001\ 0001\ 0010\ 0110\ 0100_2$

2.10(b) $8192 = 1000\ 0001\ 1001\ 0010_2$

2.11(a) Excess-3 Code

$$\begin{array}{ccc}
 6 & 7 & 8_{10} \\
 \downarrow & \downarrow & \downarrow \\
 1001 & 1010 & 1011
 \end{array}$$

2.11(b)

$$\begin{array}{ccccc}
 3 & 2 & 8 & 7 & 4_{10} \\
 \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\
 0110 & 0101 & 1011 & 1010 & 0111
 \end{array}$$

2.11(c)

$$\begin{array}{ccccc}
 6 & 1 & 4 & 4 & 0_{10} \\
 \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\
 1001 & 0100 & 0111 & 0111 & 0011
 \end{array}$$

2.12 Octal $1543 = 1 \times 8^3 + 5 \times 8^2 + 4 \times 8^1 + 3 \times 8^0$
 $= 512 + 320 + 32 + 3$
 $= \begin{array}{ccc} 8 & 6 & 7 \end{array}$
 $\downarrow \quad \downarrow \quad \downarrow$
 Excess-3 Code $\rightarrow 1011\ 1001\ 1010_2$

2.13(a) $0001\ 1001\ 0101\ 0001_2 = 4096 + 2048 + 256 + 64 + 16 + 1 = 6481 = 0110\ 0100\ 1000\ 0001\ \text{BCD}$

2.13(b) $0110\ 0001\ 0100\ 0100\ 0000_2 = 262144 + 131072 + 4096 + 1024 + 64 = 398400$
 $398400 = 0011\ 1001\ 1000\ 0100\ 0000\ 0000\ \text{BCD}$

2.14(a)

$$\begin{array}{cccccccc}
 & 1024 & 256 & 128 & & 16 & 4 & 2 & 1 & \leftarrow \text{weighing} \\
 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1_2 \\
 = & 1024 & + & 256 & + & 128 & + & 16 & + & 4 & + & 2 & + & 1 \\
 = & 1 & & 4 & & 3 & & & & & & & 1_{10} \\
 \text{Excess-3} = & 0100 & 0111 & 0110 & 0100_2
 \end{array}$$

2.14(b)

$$\begin{array}{cccccccc}
 & 1024 & 512 & & 128 & & 16 & \leftarrow \text{weighing} \\
 0 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0_2 \\
 = & 1024 & + & 512 & + & 128 & + & 16 \\
 = & 1 & 6 & 8 & 0_{10}
 \end{array}$$

$$\text{Excess-3} = 0100 \ 1001 \ 1011 \ 0011_2$$

$$2.15 \quad 1011.01$$

$$+ \quad \underline{0110.01 \ 1}$$

$$1 \ 0001.10 \ 1_2$$

$$= 16 + 1 + 0.5 + 0.125$$

$$= 17.625_{10}$$

$$2.16(a) \quad 14 = 0000 \ 1110$$

$$\underline{+17} = \underline{0001 \ 0001}$$

$$3110 = 0001 \ 1111_2$$

$$2.16(b) \quad 34 = 0010 \ 0010$$

$$\underline{+28} = \underline{0001 \ 1100}$$

$$62_{10} = 0011 \ 1110_2$$

$$2.16(c) \quad 32 = 0010 \ 0000$$

$$14 = 0000 \ 1110_2$$

$$2\text{'s Complement of } 14 = 1111 \ 0010_2$$

$$32_{10} - 14_{10} = 0010 \ 0000$$

$$\underline{+ \ 1111 \ 0010}$$

$$\text{ignore} \rightarrow 1 \ 0001 \ 0010 = 18_{10}$$

$$2.16(d) \quad 34 = 0010 \ 0010$$

$$42 = 0010 \ 1010_2$$

$$2\text{'s Complement of } 42 = 1101 \ 0110_2$$

$$34_{10} - 42_{10} = 0010 \ 0010$$

$$\underline{+ \ 1101 \ 0110}$$

$$\text{No Carry} \rightarrow 1111 \ 1000$$

$$\text{Result} = - (2\text{'s complement of } 1111 \ 1000_2)$$

$$= - (0000 \ 1000_2) = -8_{10}$$

$$2.17 \quad 3AFA_{16} = 0011 \ 1010 \ 1111 \ 1010_2$$

$$2F1E_{16} = 0010 \ 1111 \ 0001 \ 1110_2$$

$$2\text{'s complement of } 2F1E_{16} = 1101 \ 0000 \ 1110 \ 0010_2$$

$$3AFA_{16} - 2F1E_{16} = 0011 \ 1010 \ 1111 \ 1010$$

$$\underline{+ \ 1101 \ 0000 \ 1110 \ 0010}$$

$$\begin{array}{cccc}
 \text{ignore} \rightarrow 1 & \underline{0000} & \underline{1011} & \underline{1101} & \underline{1100} \\
 & 0 & B & D & C_{16}
 \end{array}$$

2.18(a) 9's complement of 132 :

$$999 - 132 = 867$$

$$\text{Hence, } 254 - 132 = 254 + 867 + 1$$

$$\begin{array}{r} = 1 \underline{122} \\ \text{ignore carry} \quad \quad \quad \uparrow \text{answer} \\ 10\text{'s complement of } 132 = 1000 - 132 = 868 \\ 254 - 132 = 254 + 868 = 1 \underline{122} \\ \text{ignore carry} \quad \quad \quad \uparrow \text{Result} = 122 \end{array}$$

2.18(b) 9's complement of 807 = 999 - 807

$$= 192$$

$$783 - 807 = 783 + 192 + 1$$

$$= 976$$

since there is no carry, result = - (1000 - 976)

$$= -24$$

10's complement of 807 = 1000 - 807

$$= 193$$

$$783 - 807 = 783 + 193 = 976$$

since there is no carry, result of subtraction = - (1000 - 976) = -24

2.19(a)

$$\begin{array}{r} 14 = 001110 \\ 8 = 001000 \\ \hline 0 010110 = +22 \\ C_5 = 0 \quad \quad \quad C_4 = 0 \\ \text{No overflow} \end{array}$$

2.19(b)

$$7 = 000111$$

$$-7 = 2\text{'s complement of } 7 = 111001$$

$$\begin{array}{r} 7 = 000111 \\ (-7) = +111001 \quad \text{no overflow} \\ \hline 0 1 000 = 0 \\ \text{Ignore} \quad \quad \quad C_5 = 1 \quad C_4 = 1 \end{array}$$

2.19 (c) 27

$$27 = 011011, \quad 19 = 010011$$

$$\begin{array}{r} +(-19) \\ \hline 8 \end{array}$$

$$-19 = 101101$$

$$\begin{array}{r} 27 = 011011 \\ +(-19) = +101101 \quad \text{no overflow} \\ \hline 8 1 000 \\ C_5 = 1 \quad \quad \quad C_4 = 1 \end{array}$$

2.19 (d) +24 = 011 000

$$-24 = 101 000$$

$$19 = 010 011$$

$$-19 = 101 101$$

$$(-24) = 101 000$$

$$+(-19) = 101 101$$

$$\begin{array}{r} -43 = 1 010 101 \quad ? \quad \text{overflow} = 1 \oplus 0 = 1 \\ C_5 = 1 \quad \quad \quad C_4 = 0 \quad \text{Hence, wrong result} \end{array}$$

2.19 (e) $19 = 010\ 011$

$$12 = 001\ 100$$

$$-12 = 110\ 100$$

$$\begin{array}{r} 19 \\ -(-12) \\ +31 \end{array}$$

$$19 = 010\ 011$$

-2's complement of -12 = +001 100

$$011\ 111 = 31$$

$$C_5 = 0 \quad C_4 = 0 \quad \text{no overflow}$$

2.19 (f) $17 = 010\ 001$

$$-17 = 101\ 111$$

$$16 = 010\ 000$$

$$-16 = 110\ 000$$

2's complement of -16 = 010 000

$$(-17) = 101\ 111$$

$$-(-16) = \underline{010\ 000}$$

$$-1 = 111\ 111 = -1_{10}$$

$$C_5 = 0 \quad C_4 = 0$$

no overflow

2.20

$$12 = 1\ 100$$

$$52 = 110\ 100$$

$$\begin{array}{r} 52 \\ \times 12 \\ \hline \end{array}$$

$$624 \quad \begin{array}{r} 1\ 100 \\ 000\ 000 \\ 0\ 000\ 00x \\ 11\ 010\ 0x \\ \hline 110\ 100\ x \end{array}$$

$$\begin{array}{r} 1\ 001\ 110\ 000 \\ \hline \end{array} = 512 + 64 + 32 + 16 = 624$$

2.21 $3 = 11_2$

$$14 = 1110_2$$

$$\frac{100}{11} = 4 = \text{Quotient}$$

$$11 \mid 111\ 0$$

$$\underline{11}$$

$$001\ 0 = 2 = \text{Remainder}$$

2.22 (a) $54 = 0101\ 0100$

$$+48 = +\underline{0100\ 1000}$$

$$102 = 1001\ 1100$$

$$\underline{0110} \quad \text{add 6}$$

$$1010\ 0010$$

$$\underline{0110} \quad \text{add 6}$$

$$0001\ 0000\ 0010 = 102 \text{ in BCD}$$

2.22 (b)

782 =	0111	1000	0010	
+219 =	0010	0001	1001	
	1001	1001	1001	1011
			0110	
	1001	1010	0001	
			0110	
	1010	0000	0001	
	0011	0000	0000	0001 = 1001 in BCD

2.22 (c)
$$\begin{array}{r} 82 = 1000\ 0010 \\ -58 = \underline{-0101\ 1000} \\ \hline 24 \end{array}$$

2's complement of each digit of 58 = 1011 1000
addition factor to find 10's complement = +1001 1010

10's complement of 58 = 1 0100 1 0010
 ignore carries

$$\begin{array}{r}
 10\text{'s complement of } 58 = 0100 \ 0010 \\
 82 = \underline{1000 \ 0010} \\
 \hline
 1100 \ 0100 \\
 \underline{0110} \\
 \hline
 \text{ignore } \nearrow 1 \quad \underbrace{0010}_2 \quad \underbrace{0100}_4
 \end{array}$$

```

2.23      999 =   1001 1001 1001
          PLUS 999 = 1001 1001 1001
          -----
          1998  1 0011 0011 0010
BCD Corrections 0110 0110 0110
                  -----
                  0001 1001 1001 1000 = 1998

```

2.24 Data with odd parity bit in MSB = 0 1011 0000
odd parity bit = 0

2.25

By conventional method (i.e. paper & pencil):

$$12 \times 52 = 624$$

Using repeated Addition :

Assume initial product to be zero. Add 52 twelve times to itself:

$$\begin{array}{l}
 0 + 52 = 52 \\
 52 + 52 = 104 \\
 104 + 52 = 156 \\
 156 + 52 = 208 \\
 208 + 52 = 260 \\
 260 + 52 = 312 \\
 312 + 52 = 364 \\
 364 + 52 = 416 \\
 416 + 52 = 468 \\
 468 + 52 = 520 \\
 520 + 52 = 572 \\
 572 + 52 = 624
 \end{array}
 \left. \vphantom{\begin{array}{l} 0 + 52 = 52 \\ 52 + 52 = 104 \\ 104 + 52 = 156 \\ 156 + 52 = 208 \\ 208 + 52 = 260 \\ 260 + 52 = 312 \\ 312 + 52 = 364 \\ 364 + 52 = 416 \\ 416 + 52 = 468 \\ 468 + 52 = 520 \\ 520 + 52 = 572 \\ 572 + 52 = 624 \end{array}} \right\} 12 \text{ times}$$

2.26

Divided
14Divisor
3

Subtraction Result

$$\begin{array}{l}
 14-3 = 11 \\
 11-3 = 8 \\
 8-3 = 5 \\
 5-3 = 2
 \end{array}$$

Counter

$$\begin{array}{l}
 1 \\
 1 + 1 = 2 \\
 2 + 1 = 3 \\
 3 + 1 = 4
 \end{array}$$

Quotient = 4, Remainder = 2

2.27

$$M = 1111\ 1111_2, Q = 1111\ 1100_2$$

Since M and Q are both negative numbers.

$$2\text{'s complement of } M = 0000\ 0001_2$$

$$2\text{'s complement of } Q = 0000\ 0100_2$$

Multiplying the 2's complement of M and Q using unsigned multiplication method,
product = $0000\ 0000\ 0000\ 0100_2 = +4_{10}$.

The sign of the product,

$$S_n = M_n \oplus Q_n = 1 \oplus 1 = 0$$

Hence, the result is $+4_{10}$

2.28

Quotient = -8, Remainder = -1. The sign of the remainder is the same as the sign of the dividend unless remainder is zero.

2.29 9-bit data with even Parity
 = 1 1100 0111 = Transmitted data has even parity.
 ↖ Parity bit
 Received data
 = 1 1000 0111 = Wrong data with odd parity.
 ↗
 Bit 6 changed from 1 to 0 .

2.30 0000 1111 = +15 in decimal
 Plus 1111 1100 = - 4 in decimal

 0000 1011 = +11 decimal

Sign bit = 0, Overflow bit = 0 since the final carry and the previous carry are the same (1's).

